



# **DATABASE - 1**

## **3<sup>RD</sup> CLASS**

### **COMPUTER SCIENCE DEPARTMENT**

5<sup>th</sup> Lecture – Mapping Relationships

Sunday 20<sup>th</sup> of Oct. 2024

**LECTURER :**

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## STEP 3 : Map Binary Relationships :

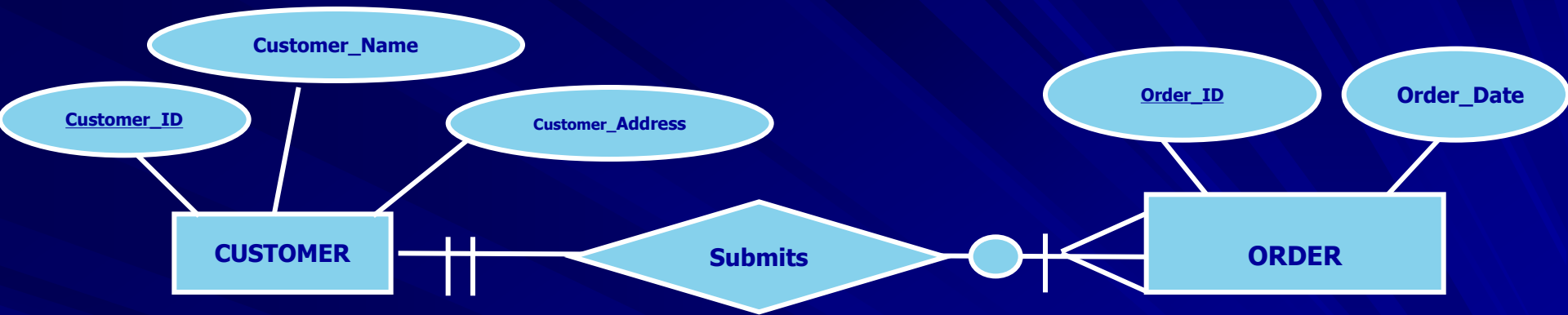
- The procedure for representing relationships depends on both
- 1. The **degree** of the relationships (**unary, binary and ternary**)
- And
- 2. The **cardinalities** of the relationships.

## A : Map Binary One-to-Many Relationships :

1. For each binary **1:M** relationship **create a relation** for each of the two entity types participating in the relationship using the procedure in Step 1.
2. **Include the primary key attribute** (or attributes) of the entity **on the one-side** of the relationship **as a foreign key** in the relation that is **one to many** – side of the relationship ( **to remember: The primary key migrates to the many side**)

We illustrate the above rules using the relationships between **CUSTOMER** and **ORDER**, which is **1:M** relationship, in the figure bellow. The primary key **Customer\_ID** of **CUSTOMER** (the one side) is included as a foreign key in **ORDER**(the many -side) .The foreign key relationship is indicated with an arrow.

Part (a)



Part (b)

**CUSTOMER**

<u>Customer Id</u>	Customer_Name	Customer_Address
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**ORDER**

<u>Order ID</u>	Order_Date	<u>Customer Id</u>
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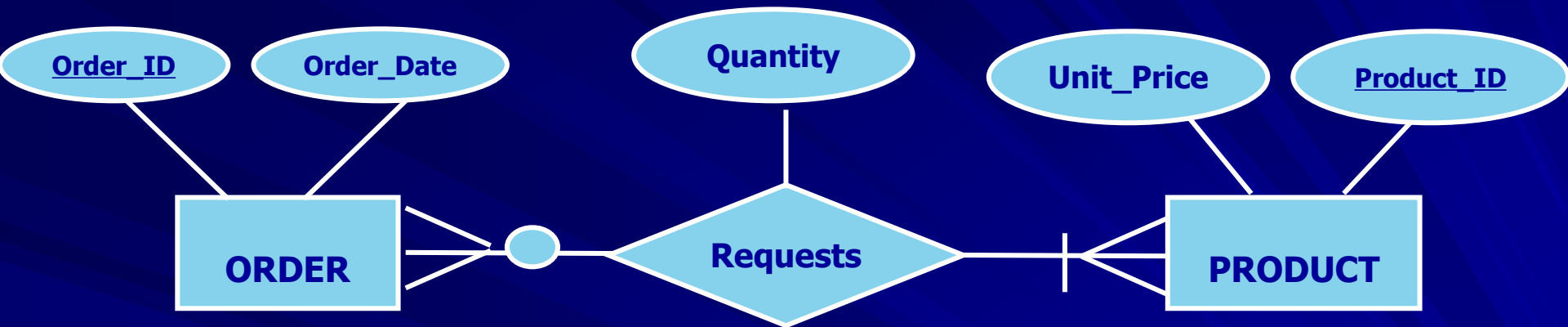


## B: Map Binary Many –to Many Relations:

- Supposing we have binary (M: N) relationship between two entity types A and B.
- For such relationship we create a new relation C. Includes as foreign key attributes in C the primary key for each of the two participating entity types.
- These attributes become the primary key of C.
- Any non-key attributes are associated with M: N relationships are included with the relation C.

- An example of applying this rule is in the figure below:
- Part (a) shows the requests relationship between entity types **ORDER** and **PRODUCT**
- Part(b) shows three relations (**ORDER**, **ORDER LINE** and **PRODUCT**) that are formed from the entity types and requests relationships.
- First, a relation created for each of the two regular entity types **ORDER**, **ORDER LINE** and **PRODUCT**.
- Then a relation (named **ORDER\_LINE** is created for the requests relationship.
- The primary key of **ORDER\_LINE** is the combination of **Order\_ID** and **Product\_ID**, which are the respective primary keys of **ORDER** and **PRODUCT** .
- As indicated in the diagram, these attributes are foreign keys that “point to “the respective keys. The non-key attribute **Quantity** also appears in **ORDER\_LINE**.

## Part (a)



## Part (b)

**ORDER**

<u>Order Id</u>	Order_Date
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**ORDER LINE**

<u>Order Id</u>	<u>Product Id</u>	Quantity
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**PRODUCT**

<u>Product Id</u>	Unit_Price	Other Attributes
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## **C: Map Binary One-to-One Relationships:**

Binary one- to -one relationships can be viewed as a special case of one- to many relationships .

The process of mapping such a relationship to relation requests two steps :

1. First, two relations are created , one for each of the participating entity types .
2. Second , the primary key of one of the relations is included as a foreign key in the other relation.

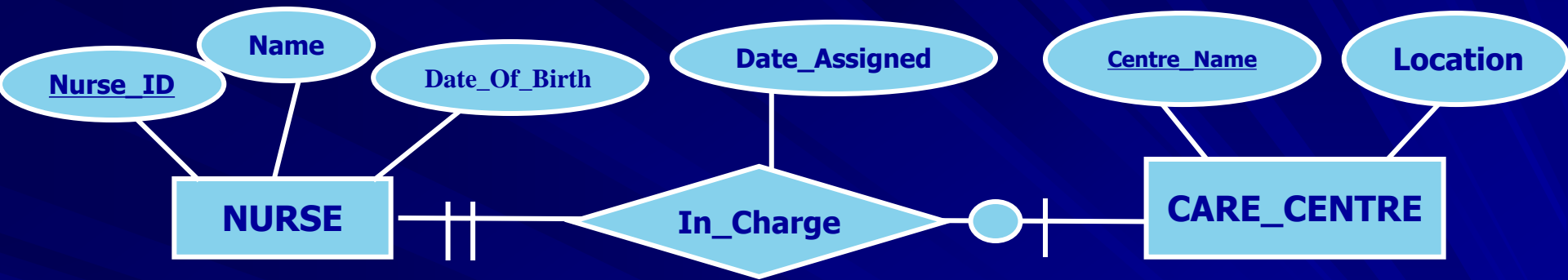


- In 1: 1 relationship, the association in one the one direction is nearly always optional one , while the association in the other direction is mandatory one .
- The foreign key of the entity type that has a mandatory participation in the 1:1 relationship should be included in the relation on the optional side of the relationship.
- This approach will avoid the need to store null values in the foreign key attribute .
- Any attributes associated with the relationship itself are also included in the same relation as the foreign key.

- As an example of applying the procedure is shown below, part (a) shows a binary 1:1 relationship between the entity type **NURSE** and **CARE\_CENTRE**.
- Each care centre must have a nurse who is in charge of that centre.
- Thus the association from **CARE\_CENTRE** to **NURSE** is a mandatory one, while the association from **NURSE** to **CARE\_CENTRE** is an optional one (since any nurse may or may not be in charge of a care centre).
- The attribute **Date\_Assigned** is attached to the **In\_Charge** relationship.

- The result of mapping this relationship to a set of relations is shown in part (b). The two relations **NURSE** and **CARE\_CENTRE** are created from the two entity types.
- Since **CARE\_CENTRE** is the optional participant, the **foreign key** is placed in this relation.
- In this case, the **foreign key** called **Nurse\_in\_Charge**. It has the same **domain** as **Nurse\_ID** , and the relationship with the primary key is shown in the figure.
- The attribute **Date\_assigned** is also located in **CARE\_CENTRE**.

Part (a)



Part (b)

**NURSE**

<u>Nurse Id</u>	Name	Date_Of_Birth
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**CARE\_CENTRE**

<u>Centre Name</u>	Location	<u>Nurse_In_Charge</u>	Date_Assigned
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## Step 4: Map Associative Entities:

- It is a good approach when the end user can best visualize the relationship as an entity type rather than as an M:N relationship.
  1. Create three relations: one for each of the participating entity types ,and the third for the associative entity. We refer to the relation formed from the associative entity as the *associative relation*.
  2. Depending on whether the on the E-R diagram an identifier was assigned to the associative entity.

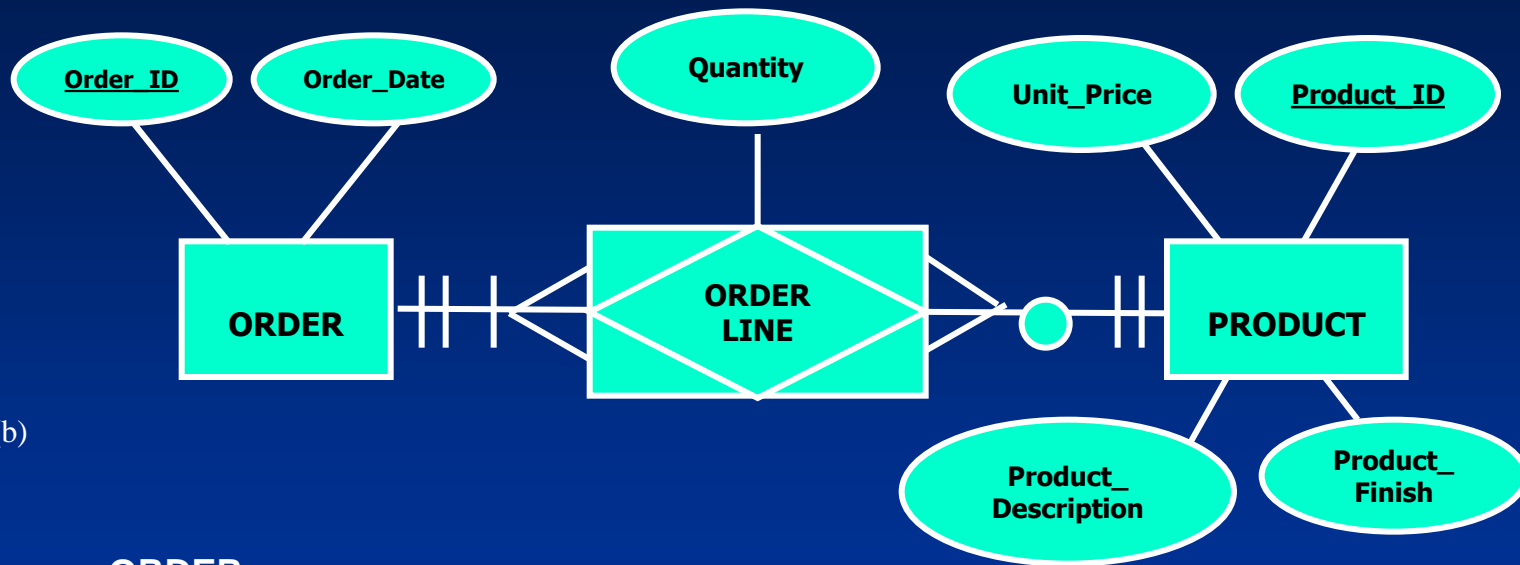


# CASE A: Identifier Not Assigned:

- If an identifier was not assigned, the default primary key for the associative relation consists of the two primary key attributes from the other two relations. These attributes are then foreign keys that reference the other two relations.
- An example of this case is shown in figure below ,part(a)shows the associative entity ORDER\_LINE that links the ORDER and PRODUCT entity types, part(b) shows the three relations that result from this mapping.



Part (a)

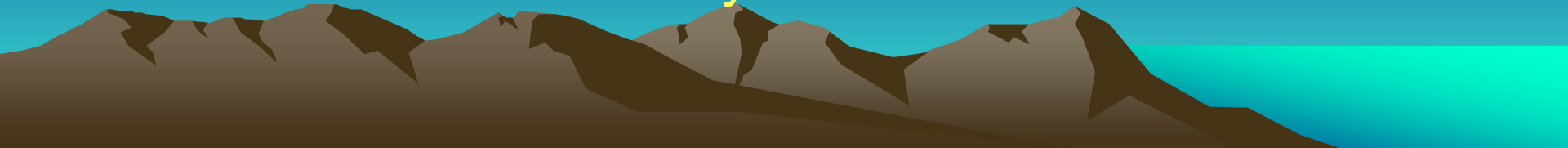


Part (b)



## CASE B: Identifier Assigned:

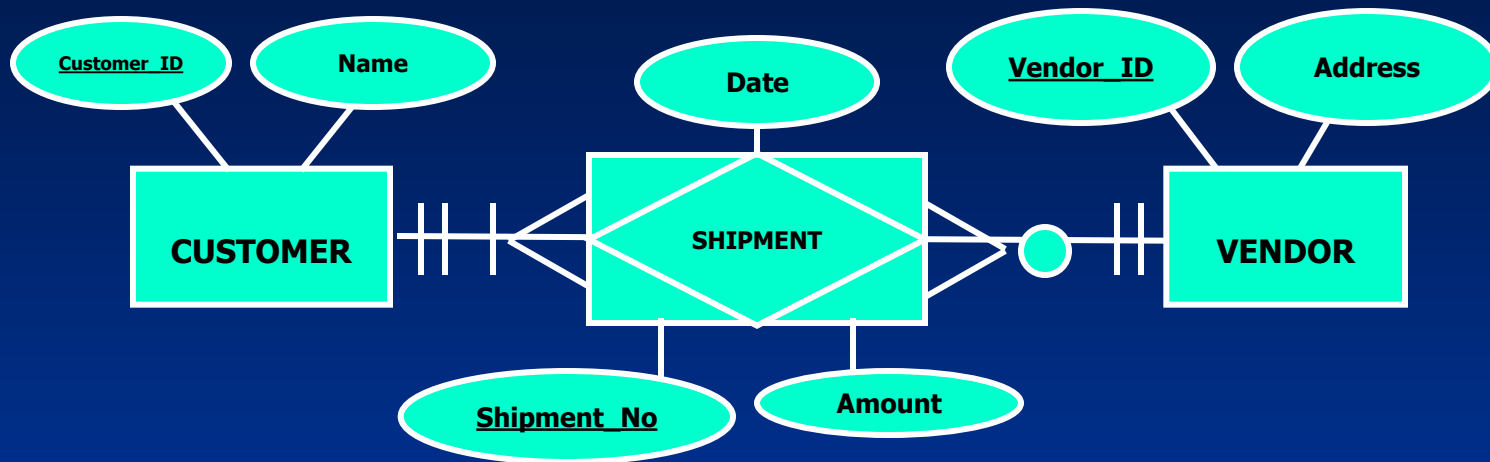
- some times the data modeler will assign an identifier (called *surrogate* identifier or key) to the associative entity type on the E-R Diagram .There are two reasons that may motivate this approach :
  1. The associative entity type has a natural identifier that is familiar to the end users.
  2. The default identifier (consisting of the identifiers for each of the participating entity types) may not uniquely identify instances of the associative entity.



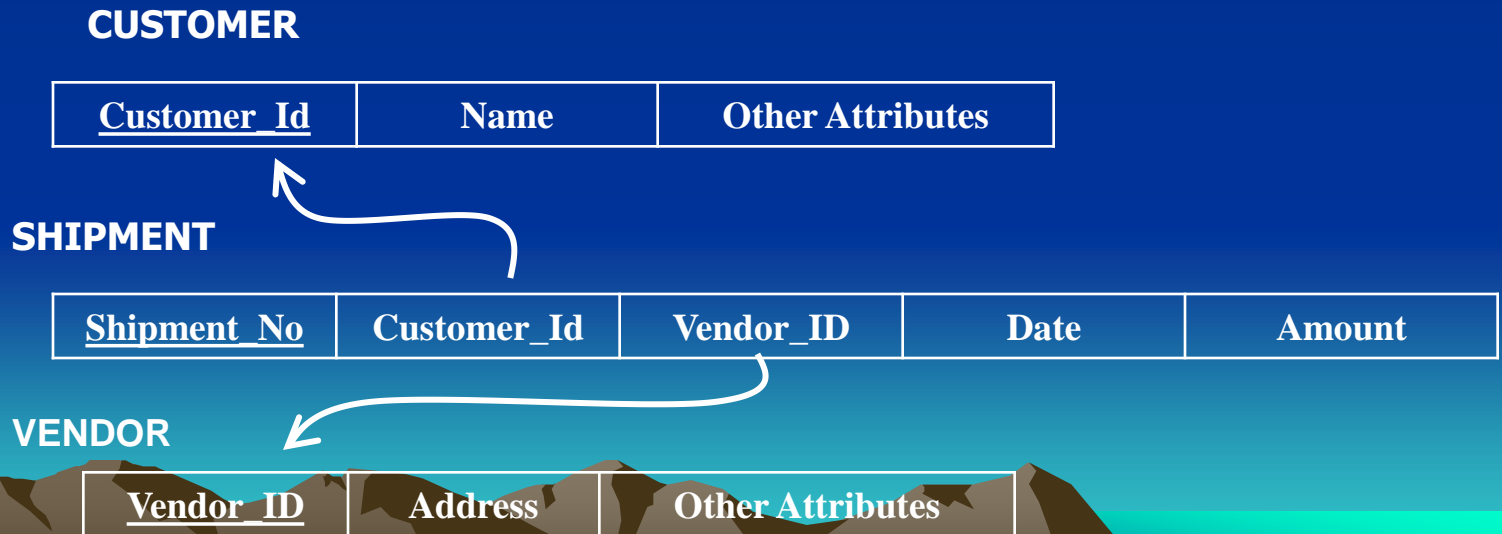


- The process for mapping the associative entity in this case is now modified as follows.
- As before, a new (associative) relation is created to represent the associative entity. However, the primary key for this relation is the identifier assigned on E-R diagram (rather than the default key ).
- The primary key for the two participating entity types are included as foreign keys in the associative relation.
- An example of this process is in the figure below, part (a) shows the associative entity type SHIPMENT that links the CUSTOMER and VENDOR entity types .Shipment\_No has been chosen as the identifier for SHIPMENT, for two reasons

Part (a)



Part (b)



Shipment\_No is a natural identifier for this entity type that is familiar to end users.

The default identifier consisting of the combination of Customer\_ID and Vendor\_Id does not uniquely identify the instances of SHIPMENT. In fact, a given vendor will make many shipments to a given customer.

Even including the attribute date does not guarantee uniqueness, since there may be more than one shipment by particular vendor on a given date, but the surrogate key Shipment\_No will uniquely identify each shipment.

Two non-key attributes associated with SHIPMENT are Date and Amount.

The result of mapping this entity to a set of relations is in part (b).

The new associative relation is named SHIPMENT.

The primary key is Shipment\_No. Customer\_ID and Vendor\_ID are included as foreign keys in this relation, and Date and Amount are Non-Key attributes.

## Step 5: Map Unary Relationships:

- Unary relationships = recursive relationships
- Two most important unary relationship cases are: one- to –many  
and many –to – many.

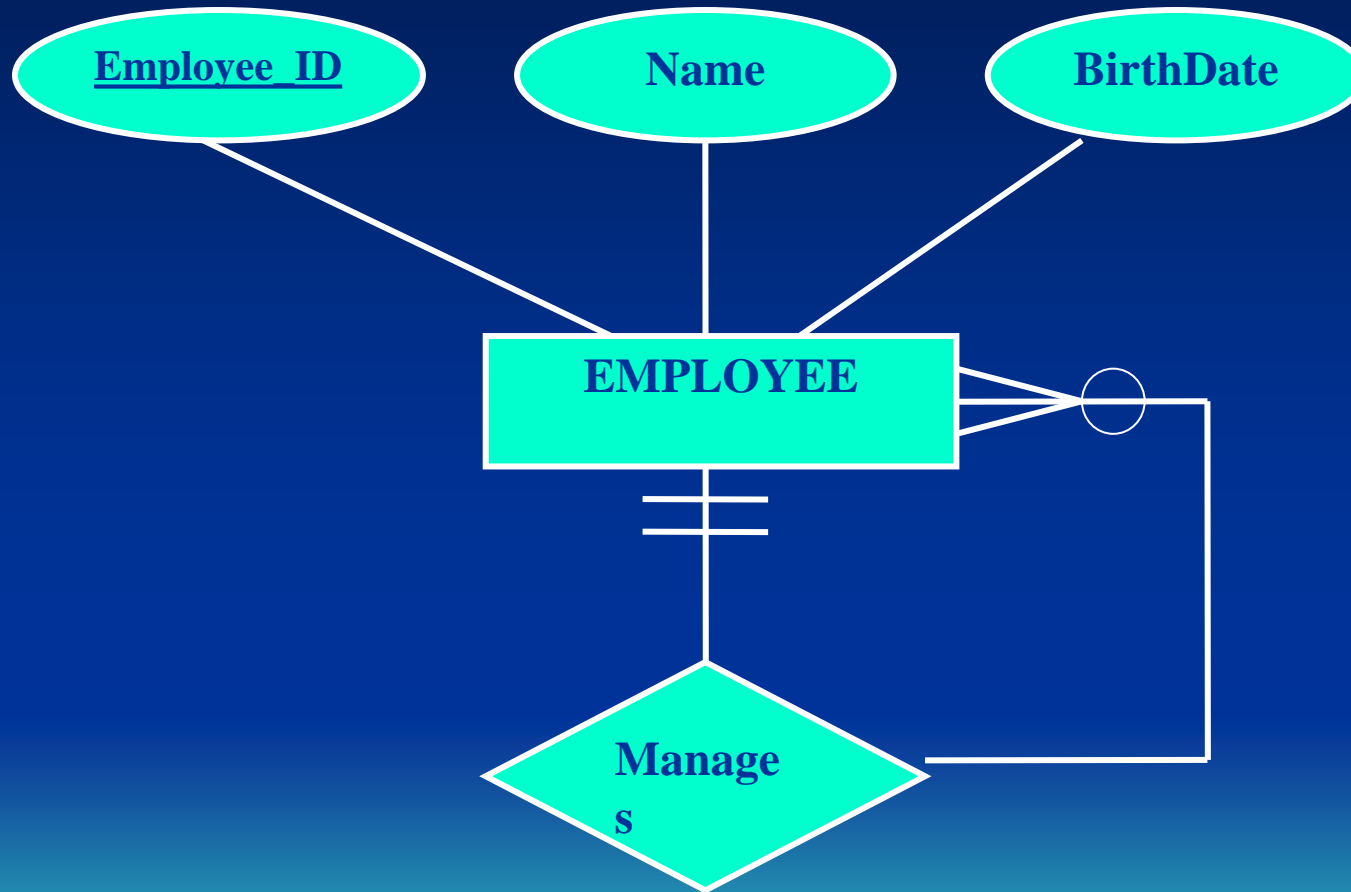




# 1. Unary one- to – many relationships:

- The entity type in the unary relationship mapped to a relation using the procedure in step (1).
- Then a foreign key attribute added *within* the same relation that references the primary key values (this foreign key must have the same domain as the primary key).
- A recursive foreign key is a foreign key in a relation that references the primary key values of the same relation .
- The figure below shows in part(a) a unary one –to-many relationship named Manages that associate each employee of an organization with another employee who is his or her manager. Each employee has exactly one manager; a given employee may manage zero to many employees .

Part (a)



- The EMPLOYEE relation that results from mapping this entity and relationship is shown in part(b) .
- The (recursive) foreign key in the relation named Manager\_ID. This attribute has the same domain as the primary key Employee\_ID.
- Each row of this relation stores the following data

Part (b)

**EMPLOYEE**

<u>Employee_ID</u>	Name	BirthDate	Manager_ID
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## 2:Unary many-to many relationships:

- In this type two relations are created: one to represent the entity type in the relation and the other an associative relation to represent the M: N relationship itself.
- The primary key of the associative relation consists of two attributes .
- These attributes (Which need not to have the same name ) both take their values from the primary key of the other relation.
- Any non-key attribute of the relationship is included in the associative relation.



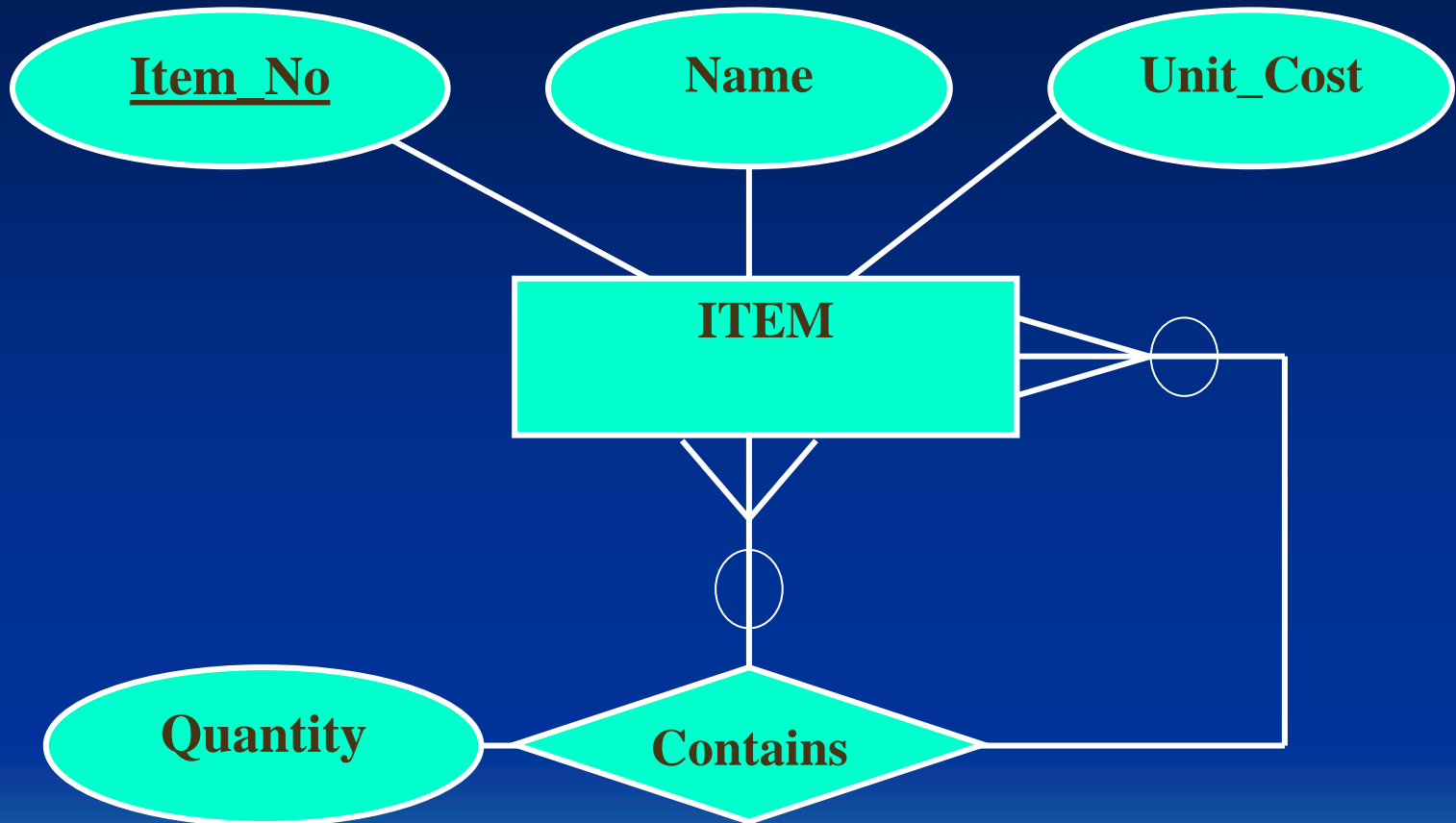


## 2:Unary many-to many relationships:

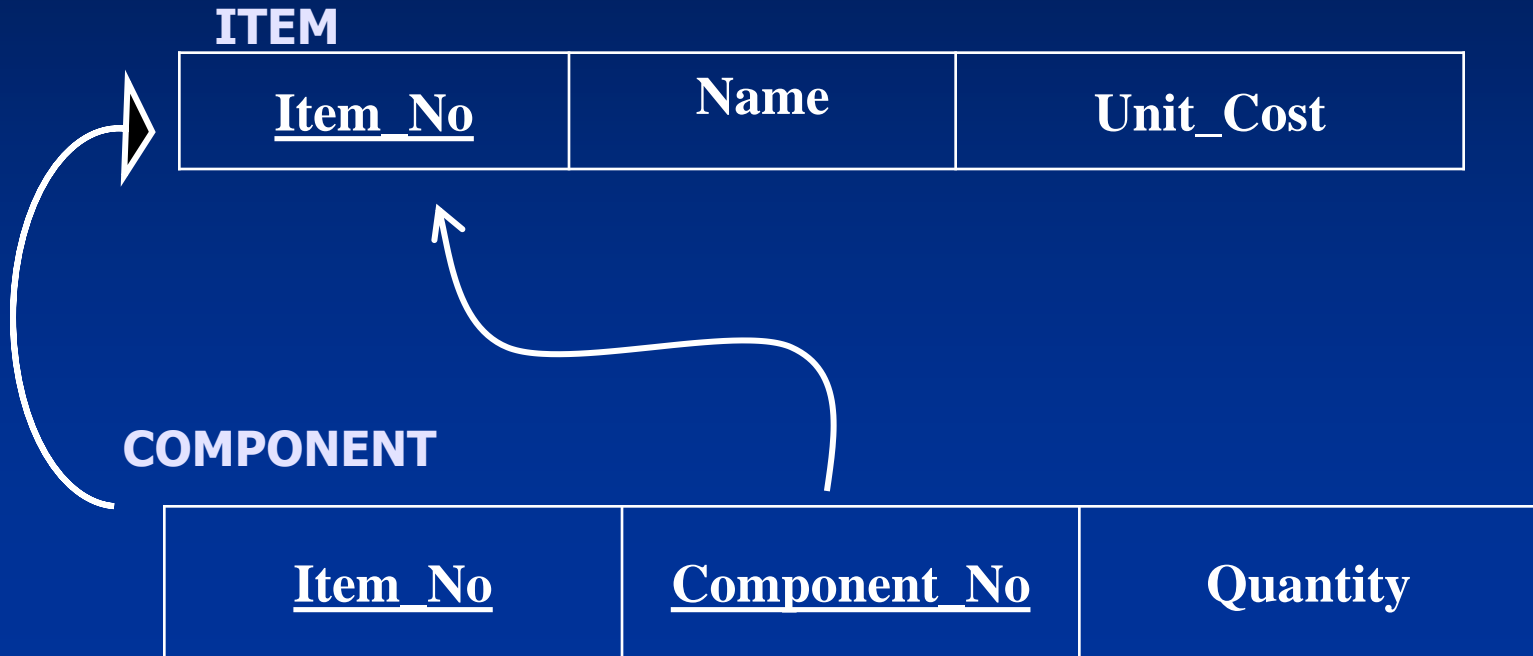
- An example of mapping a unary M:N relationship is shown in figure below, part(a) shows a bill -materials among items that are assembled from other items or components, the relationship (called Contains) is M:N since a given item can contain numerous component items, and conversely an item can be used as a component in numerous other items.



Part (a)



Part (b)



## Step 6: Map Ternary (and $n$ -ary) Relationships:

**Recall that a ternary relationship is a relationship among three entity types.**

**In fact it is recommended to convert a ternary relationship to an associative entity in order to represent participation constraints more accurately.**



## Step 6: Map Ternary (and $n$ -ary) Relationships:

**To map associative entity type that links three regular entity types or more we create a new associative relation the default primary key of this relation consists of the three primary key attributes for the participating entity types (in some cases additional attributes are required to form a unique primary key) .**

**These attributes then act in the role of foreign keys that reference the individual primary keys of the participating entity types any attributes of the associative entity type become attributes of the new relation .**

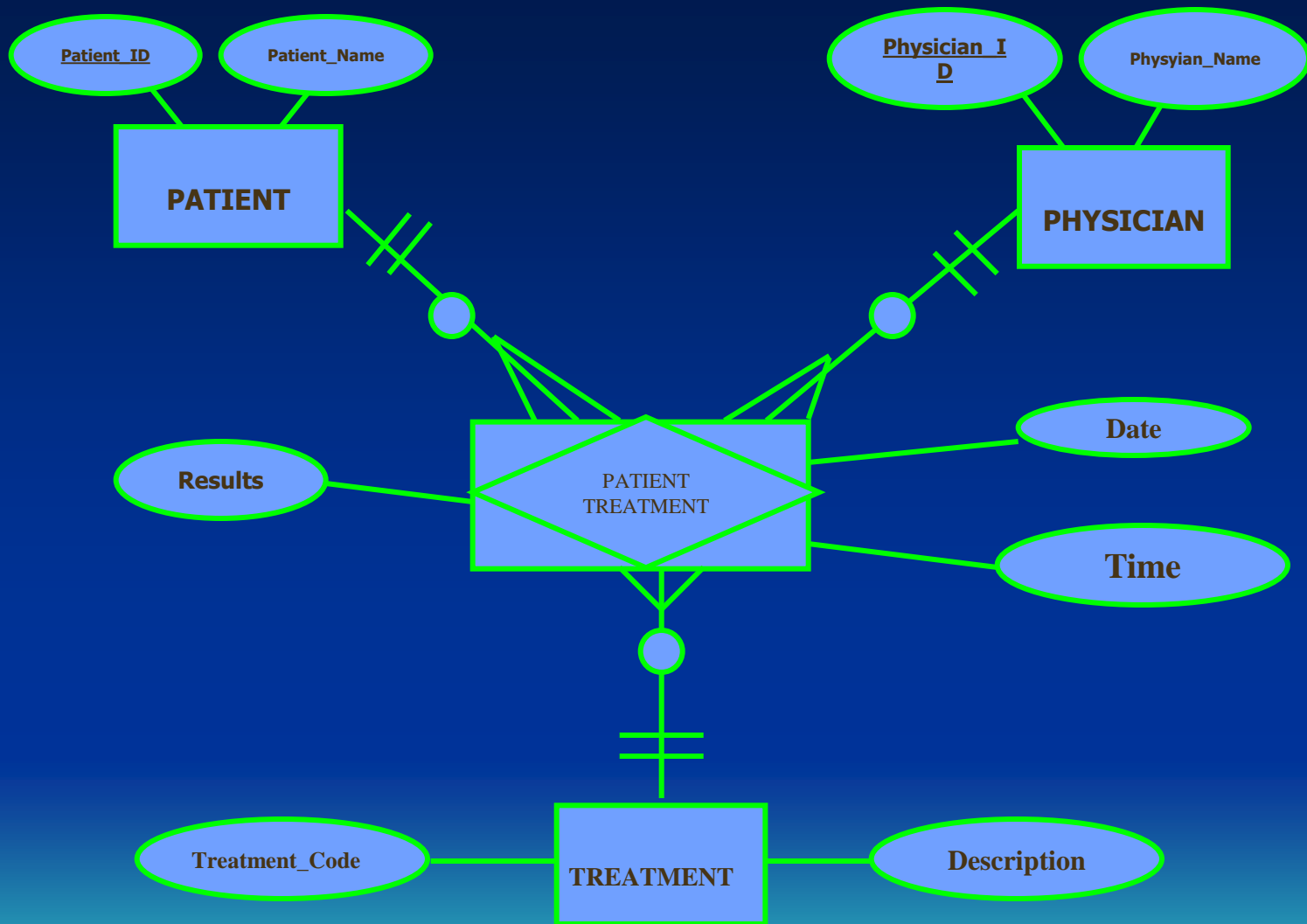


An example of mapping ternary relationship (represented as an associative entity type) is shown in figure below, in part(a) is an E-R segment (or view ) that represents *patient* receiving a *treatment* from a *physician* the associative entity type PATIENT TREATMENT has the attributes Date, Time, and results ; values are recorded for these attributes for each instance of PATIENT TREATMENT .






Part (a)



- The result of mapping this view is shown in part (b) .
- The primary key attributes Patient\_ID ,Physician\_ID,and Treatment\_Code become foreign keys in PATIENT TREATMENT. These attributes are components of the primary key of PATIENT TREATMENT.
- However they do not uniquely identify a given treatment since a patient may receive the same treatment from the same physician on more than one occasion .
- Does including the attribute Date as part of the primary key (along with the other three attributes) result in primary key ?
- This would be so if a given patient receives only one treatment from a particular physician on a given date.



- However, this is not likely to be the case. For example, a patient may be receiving a treatment in the morning , then the same treatment in the afternoon .
  - To resolve this issue we include time as apart of the primary key.
  - Therefore ,a primary key of PATIENT TREATMENT consists of the five attributes shown in part(b) :  
Patient\_UD,Physician\_ID,  
Treatment\_Code, Date,and Time.
  - The only non key attribute in the relation is Results.
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Part (b)

## PATIENT

<u>Patient_Id</u>	<u>Patient_Name</u>
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## PHYSICIAN

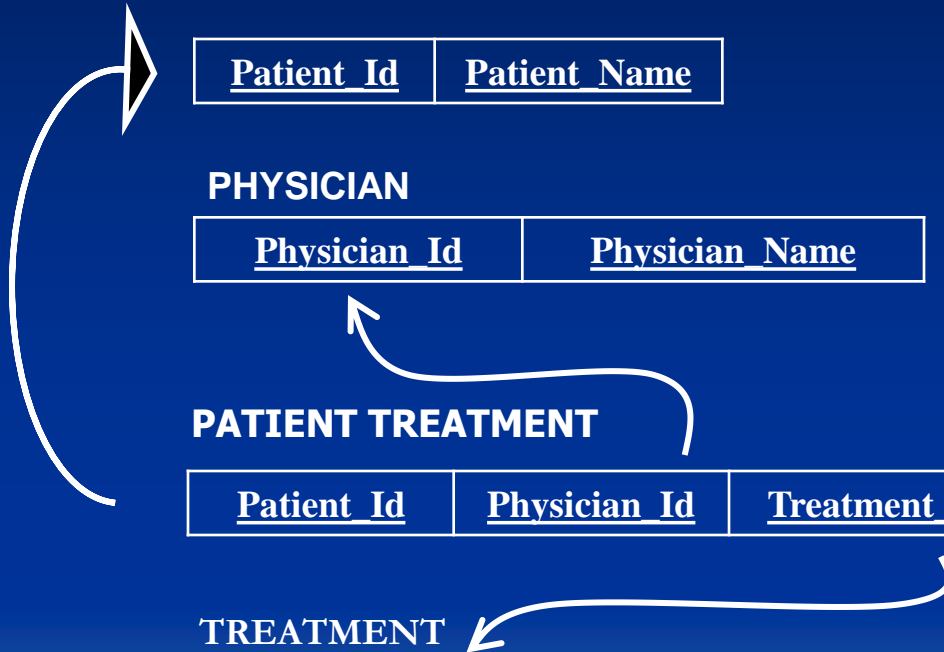
<u>Physician_Id</u>	<u>Physician_Name</u>
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## PATIENT TREATMENT

<u>Patient_Id</u>	<u>Physician_Id</u>	<u>Treatment_Code</u>	Date	Time	Results
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## TREATMENT

<u>Treatment_Code</u>	Description
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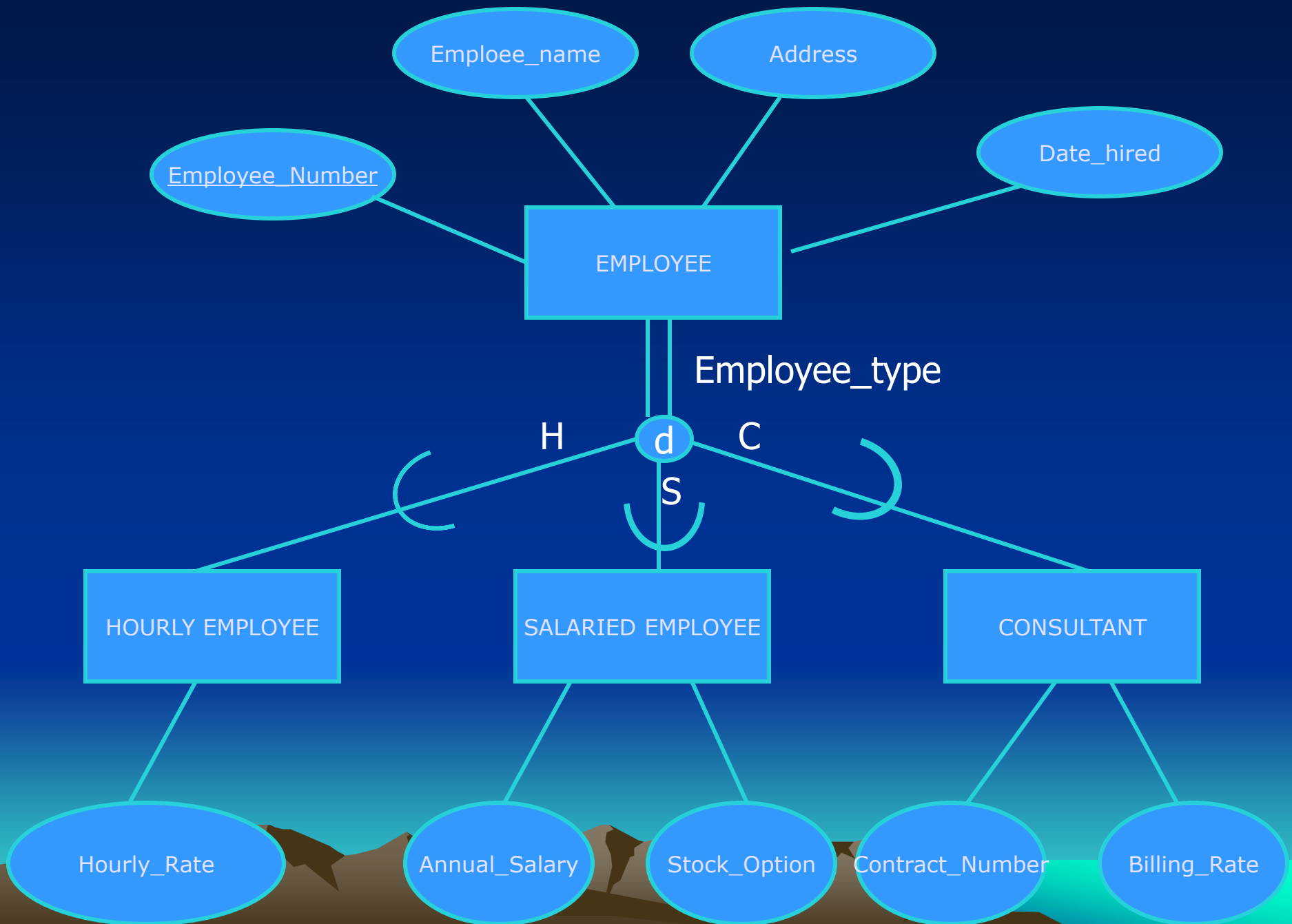
## STEP 7: MAP SUPERTYPE/SUBTYPE RELATIONSHIPS

- The relational data model does not yet directly support supertype / subtype relationships. Fortunately; there are various strategies that database designers can use to represent these relationships with the relational data model , for our purposes we use the following most commonly used strategy :
  1. Create a separate relation for the supertype and each of its subtypes .
  2. Assign to the relation created for the supertype the attributes that are commonly to all members of the supertype, including the primary key.
  3. Assign to the relation for each subtype the primary key of the supertype ,and only these attributes that are unique to that subtype.
  4. Assign one (or more) attributes of the supertype to function as the subtype discriminator.

- An example of mapping this procedure is shown in the following two figures A and B
- Figure A shows the supertype EMPLOYEE with subtypes HOURLY EMPLOYEE, SALARED EMPLOYEE, and CONSOLTANT .
- The primary key of EMPLOYEE is Employee\_Number ,and the attribute Employee\_Type is the subtype discriminator







- The result of mapping this diagram to relation using the above rule is shown in figure B.
- There is one relation for the supertype (EMPLOYEE) and one for each of the three subtypes .
- The primary key for each of the four relations is Employee\_Number.
- A prefix is used to distinguish the name of the primary key for each subtype .
- For example ,S\_Employee\_Number is the name for primary key of the relation SALARED EMPLOYEE. Each of these attributes is a foreign key that references the supertype primary key, as indicated by arrows in the diagram.
- Each subtype relation contains only those attributes peculiar to the subtype.

## EMPLOYEE

<u>Employee Number</u>	Employee_Name	Address	Employee_Type	Date_Hired
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### HOURLY\_EMPLOYEE

<u>H Employee Number</u>	Hourly_Rate
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### SALARED\_EMPLOYEE

<u>S Employee Number</u>	Annual_Salary	Stock_Options
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### CONSOLTANT

<u>C Employee Number</u>	Contract_Number	Billing_Rate
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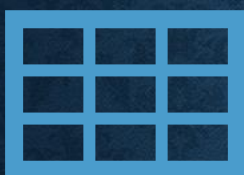


- For each subtype, a relation can be produced that contains *all* the attributes of that subtype(both specific and inherited )by using SQL command that joins the subtype with its supertype .
- For Example ,suppose that we want to display a table that contains all of the attributes for SALARED EMPLOYEE . The following command is used :

```
SELECT *  
FROM EMPLOYEE, SALARED EMPLOYEE  
WHERE Employee_Number = S_Employee_Number;
```

Although you have not yet formally studied such commands, you can intuitively see that this command will join two tables and produce a large table that contains all the attributes from both tables.





**END OF LECTURE 5**

